

**STRATEGIC BEHAVIOR UNDER COMPLETE IGNORANCE:
APPROVAL AND CONDORCET-TYPE VOTING RULES**

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ABSTRACT. Usually strategic misrepresentation of preferences in order to manipulate social choice functions is studied under the standard common knowledge assumption. In this paper, we introduce the completely opposite hypothesis of manipulation under complete ignorance. Our goal is to give an answer to the following question : do there still exist any strategic voting opportunities, even if individuals do not have any information about others' preferences ? We provide an exhaustive answer for Condorcet-type and a class of approval voting type SCFs.

1. INTRODUCTION

During the last three decades, a large literature has been devoted to the manipulation of social choice mechanisms. Gibbard (1973) and Satterthwaite (1975) proved that there does not exist any social choice procedure selecting a single alternative, which is simultaneously non dictatorial and immune to manipulation. Following these two authors, many contributors have been interested in checking the robustness of this result in different directions. For example, the equilibrium concept used by Gibbard and Satterthwaite has been relaxed in a variety of ways (see for example Pattanaik 1976, Sengupta 1978, or Mbih 1995 among others) ; another direction has been the computation of the frequency of manipulation opportunities (e.g. Lepelley and Mbih 1994). There is however one line of enquiry which has received very little attention; it concerns the possibility for individuals participating to the collective choice procedure not to be completely informed on preferences expressed by other individuals. A notable exception is due to Sengupta (1980) ; supposing a set of three possible alternatives, he shows, for a wide class of procedures based on pairwise comparisons and satisfying some very attractive properties, that in order to manipulate it may be sufficient for some individual to only know either the most preferred or the least preferred alternative of each of the other individuals.

Now what does the phrase “not completely informed” mean ? At least two interpretations are possible. First, one can think of situations, before voting takes place in a committee for instance, where some individual i knows that some other individual j prefers alternative x to alternative y , but she does not know anything else about j 's preferences, namely how j ranks x or y vis--vis any other alternative. The second interpretation is based on a probabilistic approach; given a voting rule, considering all configurations of preferences and assuming that every individual only knows his own preferences, it is possible to compute the frequency of election of each alternative; then every individual can interpret these frequencies as a distribution of probabilities over the set of alternatives;

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more concretely, before every real election, polls are generally intended to give the percentage of votes expected by every candidate, and it is theoretically possible, from the polls data, to evaluate the probability of each candidate of being elected. Voters can then use this information to choose the preferences they are going to express.

In this paper, we shall not be concerned with the probabilistic approach; we refer the reader to Andjiga, Mbih and Moyouwou (2003) for an introduction. We shall limit our analysis to the first interpretation of incomplete information, and more precisely, we focus on the extreme case in which no individual knows anything about other individuals' rankings of alternatives. Then a rational individual chooses a preference relation that better serves her interests, that is a preference that permits her in all circumstances to secure an outcome at least as good as the outcome secured by her sincere preference relation.

We examine the manipulability of the classes of voting rules. The first one contains special versions of approval voting in which a fixed number of alternatives are "approved of" by the voters ; approval voting is often said to be one of the least manipulable voting rules (see Brams and Fishburn 1984, or Yunfeng, Yue and Chen 1996) ; our results show that for the special versions we study, strategic behavior is still possible under complete ignorance, though under very binding constraints. The second set of rules we are interested in is a class of voting procedures based on the Condorcet principle : an alternative beating every other alternative in pairwise majority contests is called a Condorcet winner ; and whenever it exists, a Condorcet winner is chosen by the voting procedure.

The paper is organized as follows : section 2 introduces notations and definitions ; section 3 is concerned with results and proofs, and section 4 concludes the paper with some general remarks.

2. NOTATIONS AND DEFINITIONS

Let $N = \{1, 2, \dots, i, \dots, n\}$ be a set of n individuals ($n \geq 2$) and $A = \{a_1, a_2, \dots, a_j, \dots, a_m\}$ be a set of m alternatives ($m \geq 3$). L will denote the set of all possible linear orders on A (i.e. complete, reflexive, antisymmetric and transitive binary relation on A) and $R = a_1 a_2 a_3 \dots$ the linear order such that a_1 is the best alternative according to relation R , a_2 the second best, a_3 the third best and so on. A *profile* is an n -tuple $R^N = (R^1, \dots, R^i, \dots, R^n)$ of individual relations, one for each individual. L^N will denote the set of all possible profiles. A *contingency* is a profile from which the preference relation of some individual i has been removed, that is an $(n - 1)$ -tuple $R^{-i} = (R^1, \dots, R^{i-1}, R^{i+1}, \dots, R^n)$; and the set of all contingencies for individual i will be denoted L^{-i} . Profile R^N can now be rewritten $R^N = (R^i, R^{-i})$.

Given a non empty subset B of A , $lex(B)$ will be the alternative a_j in B with the least index j .

Definition 2.1. *An SCF f , a mapping from L^N into A , is manipulable under complete ignorance if there exists some individual i , with sincere preference R^i and strategic preference Q^i such that :*

- (a) for all $H^{-i} \in L^{-i}$, $f(Q^i, H^{-i}) R^i f(R^i, H^{-i})$
- (b) and for some $H^{-i} \in L^{-i}$, $f(Q^i, H^{-i}) \neq f(R^i, H^{-i})$

When (a) and (b) in Definition 2.1 hold, individual i has incentive to submit the false preference Q^i to the social choice procedure f , rather than his sincere preference R^i , no matter the contingency expressed by other individuals.

For example, let f be the SCF which for every profile selects the less preferred alternative according to individual 1's preference relation. Then f is manipulable under

complete ignorance : as compared with his sincere strategy, it will always be advantageous to individual 1 to express a false ordering, with her most preferred alternative ranked last.

3. RESULTS ON MANIPULATION UNDER COMPLETE IGNORANCE

3.1. Approval voting when individuals select a fixed number of alternatives.

According to the general approval principle (see Brams and Fishburn, 1983), each individual has to select one or more alternatives she most prefers, and given a tie-breaking mechanism, the social outcome is chosen from the set of the most selected alternatives. The reader can verify that such a social choice procedure is not an *SCF*; one way to define an *SCF* consists in fixing a number k ($1 \leq k \leq m - 1$) of alternatives each individual has to select.

Let $n_k(a_j, R^N)$ be the number of individuals who rank a_j among their k most preferred alternatives, and define $AP_k(R^N) = \{x \in A : n_k(x, R^N) \geq n_k(y, R^N) \text{ for all } y \in A\}$. Note that $AP_k(R^N) \neq \emptyset$ for all R^N .

Definition 3.1. An *SCF* f is a k -approval voting rule (AV_k rule) if for all profile R^N , $f(R^N) \in AP_k(R^N)$.

AV_1 and AV_{m-1} rules are versions of the well-known plurality and antiplurality rules respectively. The AV_k rule with ties broken in favor of *lex* [$AP_k(R^N)$] is called *lexicographic* AV_k rule and is denoted LAV_k .

Proposition 3.1. 1) No AV_1 rule is manipulable under complete ignorance if and only if $n \geq 3$.

2) LAV_1 is not manipulable under complete ignorance.

Proof. 1) Let f denote any AV_1 , and assume x, y and z are three distinct alternatives. Consider some individual i and let i 's sincere preference relation be $R^i = x...$, and let $Q^i = y...$ be some other preference relation for i . In order to prove that LAV_1 is manipulable under complete ignorance for $n \geq 3$, it will be sufficient to show that there exists some contingency H^{-i} such that $f(R^i, H^{-i}) \neq f(Q^i, H^{-i})$ and $f(R^i, H^{-i}) R^i f(Q^i, H^{-i})$. We shall distinguish two cases :

(a) $n \geq 3$ and odd. Consider two disjoint subsets N_1 and N_2 of $N - \{i\}$ with $\frac{n-1}{2}$ members each and construct contingency H^{-i} such that for all $i \in N_1$, $H^i = xy...$ and for all $i \in N_2$, $H^i = yx...$.

(b) $n \geq 4$ and even. Consider some $j \neq i$ and two disjoint subsets N_1 and N_2 of $N - \{i, j\}$ with $\frac{n}{2} - 1$ members each, and construct contingency H^{-i} such that $H^j = z...$, for all $i \in N_1$, $H^i = xy...$ and for all $i \in N_2$, $H^i = yx...$.

Then, in both cases, for all AV_1 's f , $f(R^i, H^{-i}) = x$, $f(Q^i, H^{-i}) = y$, and xR^iy . Therefore no AV_1 is manipulable under complete ignorance. It remains to show that for $n = 2$, there exists some AV_1 which is manipulable under complete ignorance. Suppose $n = 2$, and let g be an AV_1 defined as follows :

$$\begin{aligned} g(R^1, R^2) &= a_j \text{ for all } a_j \text{ if } a_2 \text{ is top in } R^1 \text{ and } a_j \text{ is top in } R^2 \\ g(R^1, R^2) &= a_2 \text{ if } a_1 \text{ is top in } R^1 \text{ and } a_2 \text{ is top in } R^2 \\ g(R^1, R^2) &= a_1 \text{ if } a_1 \text{ is top in } R^1 \text{ and } a_2 \text{ is not top in } R^2 \\ \text{and } g(R^1, R^2) &= LAV_1(R^1, R^2) \text{ if } a_j \text{ is top in } R^1 \text{ and } a_j \notin \{a_1, a_2\}. \end{aligned}$$

Clearly g is manipulable under complete ignorance, by replacing a sincere preference $R^i = a_2a_1...$ with some other preference $Q^i = a_1...$.

2) From 1), if $n \geq 3$, LAV_1 is not manipulable under complete ignorance. For $n = 2$, consider some individual i ; it will be sufficient, as above, to show that there exists some contingency H^j such that $f(R^i, H^j) \neq f(Q^i, H^j)$ and $f(R^i, H^j) R^i f(Q^i, H^j)$. Let $R^i = a_k \dots$ and $Q^i = a_h \dots$. Choose $H^j = a_k \dots$ if $h < k$ and $H^j = a_h \dots$ if $h > k$. It clearly follows that LAV_1 is not manipulable under complete ignorance. \square

Proposition 3.2. *Assume $2 \leq k \leq m - 1$.*

1) *LAV_k is manipulable under complete ignorance if $k > \frac{1+m(n-1)}{n}$.*

2) *No AV_k is manipulable under complete ignorance if and only if $k \leq \frac{1+m(n-1)}{n}$.*

Proof : 1) Suppose $k > \frac{1+m(n-1)}{n}$. Without loss of generality suppose individual 1's sincere preference relation is $R^1 = a_2 a_3 \dots a_k a_1 a_{k+1} \dots a_m$ and consider $Q^1 = a_2 a_3 \dots a_k a_{k+1} \dots a_m a_1$. We shall show that LAV_k is manipulable under complete ignorance by submitting Q^1 instead of R^1 .

Let $A_k(R)$ be the set of the k first most preferred alternatives according to linear order R . For any contingency H^{-1} , we then can write

$$A = \left[\bigcap_{i \neq 1} A_k(H^i) \right] \cup \left[\bigcup_{i \neq 1} [A - A_k(H^i)] \right].$$

Clearly $\bigcap_{i \neq 1} A_k(H^i)$ and $\bigcup_{i \neq 1} [A - A_k(H^i)]$ are disjoint, hence

$$\text{Card} \left[\bigcap_{i \neq 1} A_k(H^i) \right] = \text{Card} A - \text{Card} \left[\bigcup_{i \neq 1} [A - A_k(H^i)] \right],$$

and therefore

$$(3.1) \quad \text{Card} \left[\bigcap_{i \neq 1} A_k(H^i) \right] \geq \text{Card} A - \sum_{i \neq 1} \text{Card} [A - A_k(H^i)].$$

For each $i \in N - \{1\}$, $\text{Card} [A - A_k(H^i)] = m - k$, then from (3.1) and the fact that $k > \frac{1+m(n-1)}{n}$ can be rewritten $k - 1 > (n - 1)(m - k)$ we must have :

$$(3.2) \quad \text{Card} \left[\bigcap_{i \neq 1} A_k(H^i) \right] \geq m - (n - 1)(m - k) > m - (k - 1)$$

Inequality (3.2) implies that

$$B = \{a_2, a_3, \dots, a_k\} \cap \left[\bigcap_{i \neq 1} A_k(H^i) \right] \neq \emptyset.$$

It follows that

$$LAV_k(Q^1, H^{-1}) = \text{lex}(B), \text{ and } LAV_k(R^1, H^{-1}) \in \{\text{lex}(B), a_1\}.$$

Thus $LAV_k(Q^1, H^{-1}) R^1 LAV_k(R^1, H^{-1})$. And moreover, for all contingencies H^{-1} such that $a_1 \in A_k(H^i)$ for every $i \in N - \{1\}$, $LAV_k(R^1, H^{-1}) = a_1$. And finally we conclude that individual 1 can manipulate LAV_k under complete ignorance.

2) Suppose $k \leq \frac{1+m(n-1)}{n}$. First of all, note that this can be rewritten

$$(3.3) \quad k - 1 \leq (n - 1)(m - k)$$

Consider individual 1's sincere preference R^1 and some other preference Q^1 . It is clear that if $A_k(R^1) = A_k(Q^1)$, then individual 1 cannot change the social outcome of an AV_k rule by submitting Q^1 instead of R^1 . Now suppose that $A_k(R^1) \neq A_k(Q^1)$. Since $A_k(R^1)$ and $A_k(Q^1)$ have the same cardinality, we must have $A_k(R^1) - A_k(Q^1) \neq \emptyset$ and $A_k(Q^1) - A_k(R^1) \neq \emptyset$. Let $x \in A_k(R^1) - A_k(Q^1)$ and $y \in A_k(Q^1) - A_k(R^1)$ and define two

integers q and r in the following way : $k - 1 = q(m - k) + r$ with $0 \leq r < m - k$. As a consequence of (3.3), we must have $q \leq n - 1$ if $r = 0$ and $q + 1 \leq n - 1$ if $r > 0$. Construct $A_k(R^1) - \{x\} = \bigcup_{2 \leq i \leq q+2} E_i$ where E_2, E_3, \dots, E_{q+2} are disjoint subsets of $A_k(R^1) - \{x\}$ such that $\text{card}E_i = m - k$ if $2 \leq i \leq q + 1$ and $\text{card}E_{q+2} = r$.

And consider some contingency H^{-1} such that :

$$(3.4) \quad \begin{cases} (a) \text{ for all } i \in N - \{1\}, \{x, y\} \subseteq A_k(H^i) \\ (b) \text{ for all } i \in \{2, 3, \dots, q + 2\}, E_i \subseteq A - A_k(H^i) \end{cases}$$

Such a contingency can be obtained when every individual in $N - \{1\}$ ranks x first and y second, and each $i = 2, 3, \dots, q + 2$ does not rank any alternative in E_i among her k most preferred alternatives. From (3.4), $AP_k(R^1, H^{-1}) = \{x\}$ and $AP_k(Q^1, H^{-1}) = \{y\}$; and therefore for all AV_k rule f , $f(R^1, H^{-1}) = x$ and $f(Q^1, H^{-1}) = y$. Since xR^1y , it follows that individual 1 cannot manipulate f under complete ignorance by submitting Q^1 instead of R^1 .

When $k > \frac{1+m(n-1)}{n}$, there exist some AV_k rules (e. g. LAV_k) which are manipulable under complete ignorance, as shown above. \square

The statements below are straightforward consequences of the results above.

Corollary 3.1. 1) No AV_2 is manipulable under complete ignorance.

2) LAV_k is manipulable under complete ignorance if and only if $k > \frac{1+m(n-1)}{n}$.

3) Suppose $2 \leq k < m - 1$. If LAV_k is manipulable under complete ignorance, then so is LAV_{k+1} .

4) LAV_{m-1} is manipulable under complete ignorance if and only if $m \geq n + 2$.

5) If $m \leq n + 1$, then no AV_k rule is manipulable under complete ignorance.

Table 1 and Table 2 summarize these results. It appears that, roughly speaking, profitable misrepresentation of preferences under complete ignorance is possible only when the number of alternatives is greater than the number of voters. This is usually the case in small committees dealing for example with multi-candidate applications, like in then case of recruitment of lecturers in universities. But with large electorates, like in political elections, this will never occur.

3.2. SCFs based on Condorcet principle. Given a profile R^N , the set $C(R^N)$ of Condorcet winners of R^N is the set of all alternatives that are preferred to any other alternative by at least half of the number of individuals.

Definition 3.2. f is a Condorcet-type SCF ($CT - SCF$) if for every profile R^N , $f(R^N) \in C(R^N)$ whenever $C(R^N)$ is a non empty subset of A .

It is well-known that $C(R^N)$ can be empty or contain more than one alternative for n even. We shall call *lexicographic CT - SCFs*, denoted $LCT - SCFs$, the $CT - SCFs$ that always break ties in favor of $\text{lex}[C(R^N)]$.

Proposition 3.3. No $CT - SCF$ is manipulable under complete ignorance if and only if n is odd.

Proof. (a) First suppose n is odd. Without loss of generality consider individual 1's sincere preference R^1 and some other preference Q^1 . Let $\{x, y\} \subseteq A$ be such that xR^1y and yQ^1x , N_1 and N_2 be two disjoint subsets of $N - \{1\}$ with $\frac{n-1}{2}$ members each, and H^{-1} be some contingency such that for all $i \in N_1$, $H^i = xy\dots$ and for all $i \in N_2$, $H^i = yx\dots$. Each $CT - SCF$ f verifies $f(R^1, H^{-1}) = x$ and $f(Q^1, H^{-1}) = y$. Thus f is not manipulable under complete ignorance by submitting Q^1 instead of R^1 .

(b) Now suppose n is even. Consider $\{x, y, z\} \subseteq A$ and f any $CT - SCF$ where ties are broken as follows :

(T1) $f(R^N) = x$ if $C(R^N) = \{x, y\}$ and $n(x, y, R^N) \geq \frac{n}{2}$ where $n(x, y, R^N) = \text{Card}\{i \in N : xR^i y\}$.

(T2) $f(R^N) = y$ if $C(R^N) = \{x, y\}$ and $n(x, y, R^N) < \frac{n}{2}$.

(T3) $f(R^N) = y$ if $y \in C(R^N)$ and $C(R^N) - \{x, y\} \neq \emptyset$.

(T4) $f(R^N) = \text{lex}[C(R^N) - \{x\}]$ if $y \notin C(R^N)$ and $C(R^N) - \{x\} \neq \emptyset$.

(T5) $f(R^N) = x$ if $C(R^N) = \emptyset$, $n(x, y, R^{-1}) = \frac{n}{2} - 1$ and $n(x, a, R^{-1}) \geq \frac{n}{2}$ for all $a \neq x$.

(T6) $f(R^N) = z$ if $C(R^N) = \emptyset$ and $(n(x, y, R^{-1}) \neq \frac{n}{2} - 1$ or $n(x, a, R^{-1}) < \frac{n}{2}$ for some $a \neq x$).

Consider individual 1's sincere preference $R^1 = xy\dots z$ and $Q^1 = yx\dots z$ such that R^1 and Q^1 differ only on $\{x, y\}$. We shall show that f is manipulable under complete ignorance by submitting Q^1 instead of R^1 .

Every contingency H^{-1} verifies $C(R^1, H^{-1}) - \{x\} \subseteq C(Q^1, H^{-1}) \subseteq C(R^1, H^{-1}) \cup \{y\}$.

Suppose $C(Q^1, H^{-1}) \neq C(R^1, H^{-1})$. One of the following cases holds.

Case 1 : $y \notin C(Q^1, H^{-1})$. Therefore $C(Q^1, H^{-1}) = C(R^1, H^{-1}) - \{x\}$ and $y \notin C(R^1, H^{-1})$.

If $C(Q^1, H^{-1}) \neq \emptyset$, then from (T4), $f(Q^1, H^{-1}) = f(R^1, H^{-1}) = \text{lex}[C(R^1, H^{-1}) - \{x\}]$.

If $C(Q^1, H^{-1}) = \emptyset$, then $C(R^1, H^{-1}) = \{x\}$ and from (T5), $f(Q^1, H^{-1}) = f(R^1, H^{-1}) = x$.

Case 2 : $y \in C(Q^1, H^{-1})$ and $C(R^1, H^{-1}) - \{x, y\} \neq \emptyset$. From (T3) and (T4), $f(Q^1, H^{-1}) = y$ and $f(R^1, H^{-1}) \neq x$. So $f(Q^1, H^{-1})R^1 f(R^1, H^{-1})$.

Case 3 : $y \in C(Q^1, H^{-1})$ and $C(R^1, H^{-1}) - \{x, y\} = \emptyset$. Therefore $C(Q^1, H^{-1}) \subseteq \{x, y\}$.

If $C(Q^1, H^{-1}) = \{x, y\}$, then $C(R^1, H^{-1}) = \{x\}$, $f(R^1, H^{-1}) = x$ and from (T1), $f(Q^1, H^{-1}) = x$.

But if $C(Q^1, H^{-1}) = \{y\}$, then $C(R^1, H^{-1}) = \{x, y\}$ or $C(R^1, H^{-1}) = \emptyset$. From (T2) and (T6), $f(Q^1, H^{-1}) = y$, and $f(R^1, H^{-1}) = y$ or $f(R^1, H^{-1}) = z$.

For each case above we obtain $f(Q^1, H^{-1})R^1 f(R^1, H^{-1})$. Moreover for some contingency H^{-1} such that $\frac{n}{2}$ individuals have preferences $xy\dots$ and $\frac{n}{2} - 1$ other individuals have preferences $yxz\dots$, we must have $C(R^1, H^{-1}) = \{x, z\}$ and $C(Q^1, H^{-1}) = \{x, y, z\}$. Then from (T3) and (T4), $f(Q^1, H^{-1}) = y$ and $f(R^1, H^{-1}) = z$ respectively. Hence f is manipulable under complete ignorance by submitting Q^1 instead of R^1 .

Proposition 3.4. *No LCT - SCF is manipulable under complete ignorance if and only if $n \neq 2$.*

Proof : From Proposition 3.3, no $LCT - SCF$ is manipulable under complete ignorance if n is odd. Now suppose that n is even and $n \geq 4$. Let $R^1, Q^1 \in L$, $\{a_j, a_k\} \subseteq A$ such that $j < k$, $a_j R^1 a_k$ and $a_k Q^1 a_j$; let N_1 and N_2 be two disjoint subsets of $N - \{1\}$ with $\frac{n}{2} - 1$ and $\frac{n}{2}$ members respectively; consider $H^{-1} \in L^{-1}$ such that for all $i \in N_1$, $H^i = a_j a_k \dots$ and for all $i \in N_2$, $H^i = a_k a_j \dots$. Since $C(R^1, H^{-1}) = \{a_j, a_k\}$ and $C(Q^1, H^{-1}) = \{a_k\}$, $f(R^1, H^{-1}) = a_j$ and $f(Q^1, H^{-1}) = a_k$ for any $LCT - SCF$. Therefore f is not manipulable under complete ignorance by submitting Q^1 instead of R^1 .

Suppose $n = 2$. The reader can check that there is a unique $LCT - SCF$ f which is manipulable under complete ignorance by submitting $Q^1 = a_2 a_3 a_1 \dots$ instead of $R^1 = a_2 a_1 a_3 \dots$ as shown in the table below.

	individual 2's ranking on $\{a_1, a_2, a_3\}$					
	$a_1 a_2 a_3$	$a_1 a_3 a_2$	$a_2 a_1 a_3$	$a_2 a_3 a_1$	$a_3 a_1 a_2$	$a_3 a_2 a_1$
$R^1 = a_2 a_1 a_3 \dots$	a_1	a_1	a_2	a_2	a_1	a_2
$Q^1 = a_2 a_3 a_1 \dots$	a_1	a_1	a_2	a_2	a_2	a_2

□

4. CONCLUDING REMARKS

In this paper, we consider the manipulability of two classes of *SCFs*, when individuals completely ignore how other agents rank alternatives. In that context, we show that Condorcet-based *SCFs* are in the general case immune to strategic manipulation when the number of individuals is odd. And for *LCT – SCFs*, there remain limited opportunities for strategic voting only when the society is reduced to two individuals. For the class of approval voting rules under consideration, the answer is less optimistic: we can find situations in which, by misrepresenting her preferences, some individual can without any risk secure an outcome she prefers to the outcome chosen under sincere preferences. Nevertheless, it is remarkable that, in contrast with the striking negative feature of the Gibbard-Satterthwaite theorem, most of the statements proved in this paper are rather positive. Furthermore, even pathological situations are clearly extreme cases, as shown in the Tables. Now, it would be interesting to extend these results into at least three other distinct - though not disjoint - lines of enquiry: 1) the study of manipulation under complete ignorance for other classes of *SCFs*, like scoring voting methods for example, 2) the comparison of *SCFs* on the basis of the minimum level of information necessary for manipulation, and 3) the study of manipulation under a probabilistic framework.

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Appendix

Table 1. Upper bound values of k such that no AV_k rule is manipulable under complete ignorance

$n \downarrow m \rightarrow$	3	4	5	15	30	45	60	75	90	105	120	135	150	165	180	195
2*	2	2	3	8	15	23	30	38	45	53	60	68	75	83	90	98
3	2	3	3	10	20	30	40	50	60	70	80	90	100	110	120	130
4	2	3	4	11	22	34	45	56	67	79	90	101	112	124	135	146
5	2	3	4	12	24	36	48	60	72	84	96	108	120	132	144	156
15	2	3	4	14	28	42	56	70	84	98	112	126	140	154	168	182
30	2	3	4	14	29	43	58	72	87	101	116	130	145	159	174	188
45	2	3	4	14	29	44	58	73	88	102	117	132	146	161	176	190
60	2	3	4	14	29	44	59	73	88	103	118	132	147	162	177	191
75	2	3	4	14	29	44	59	74	88	103	118	133	148	162	177	192
90	2	3	4	14	29	44	59	74	89	103	118	133	148	163	178	192
105	2	3	4	14	29	44	59	74	89	104	118	133	148	163	178	193
120	2	3	4	14	29	44	59	74	89	104	119	133	148	163	178	193
135	2	3	4	14	29	44	59	74	89	104	119	134	148	163	178	193
150	2	3	4	14	29	44	59	74	89	104	119	134	149	163	178	193
165	2	3	4	14	29	44	59	74	89	104	119	134	149	164	178	193
180	2	3	4	14	29	44	59	74	89	104	119	134	149	164	179	193
195	2	3	4	14	29	44	59	74	89	104	119	134	149	164	179	194

*For $n=2$, k must be different from 1.

Table 2. Lower bound values of n such that no AV_k rule is manipulable under complete ignorance

$k \downarrow m \rightarrow$	3	4	5	6	7	8	9	10	11	12	13
2	2	2	2	2	2	2	2	2	2	2	2
3	—	3	2	2	2	2	2	2	2	2	2
4	—	—	4	3	2	2	2	2	2	2	2
5	—	—	—	5	3	3	2	2	2	2	2
6	—	—	—	—	6	4	3	3	2	2	2
7	—	—	—	—	—	7	4	3	3	3	2
8	—	—	—	—	—	—	8	5	4	3	3
9	—	—	—	—	—	—	—	9	5	4	3
10	—	—	—	—	—	—	—	—	10	6	4
11	—	—	—	—	—	—	—	—	—	11	6
12	—	—	—	—	—	—	—	—	—	—	12

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